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STRUCTURAL PROPERTIES OF ENDOFUNCTORS

by A. BARKHUDARYAN, V. KOUBEK AND V. TRNKOVA

ABSTRACT. Un foncteur $F : \mathbb{K} \rightarrow \mathbb{L}$ est un DVO-foncteur s'il est naturellement équivalent à tout foncteur $G : \mathbb{K} \rightarrow \mathbb{L}$ tel que pour tout \mathbb{K} -object X , FX soit isomorphe à GX . On démontre que chaque DVO-foncteur $F : \text{SET} \rightarrow \text{SET}$ est finitaire (c.-à-d., préserve les colimites dirigées).

1. INTRODUCTION AND MAIN THEOREM

Inspired by [6,7], systems of functorial equations were introduced and investigated in [10]. These are systems of equations of the form

$$\mathbb{F}(\alpha) = \beta$$

where \mathbb{F} is a functorial symbol and α, β are cardinal numbers. A functor $F : \text{SET} \rightarrow \text{SET}$ is a solution of a system \mathcal{S} if, for every equation $\mathbb{F}(\alpha) = \beta$ of \mathcal{S} ,

$$\text{card } F(\alpha) = \beta.$$

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Clearly, if F is a solution of \mathcal{S} , then every functor naturally equivalent to F is a solution of \mathcal{S} as well.

Following [10], we say that a system \mathcal{S} of functorial equations is solvable (or uniquely solvable) if it has a solution (or a solution unique up to natural equivalence).

In [10], the solvability of the systems of two functorial equations

$$\mathbb{F}(\alpha_1) = \beta_1$$

$$\mathbb{F}(\alpha_2) = \beta_2$$

is discussed in the dependence of the quadruple $(\alpha_1, \alpha_2, \beta_1, \beta_2)$ of cardinal numbers. In ‘almost all’ cases, the decision whether the system is solvable or not is presented in [10]. For the cases remaining open in [10], it is impossible to give a simple YES/NO answer to the question about the solvability of the system because, as proved in [4], the answer depends on the set-theory used. In contrast to this, the following statement is absolute:

the solution of an arbitrary uniquely solvable system of functorial equations is a finitary functor (i.e., one which preserves directed colimits).

In fact, every functor $F : \mathbf{SET} \rightarrow \mathbf{SET}$ determines its canonical system of functorial equations, namely the system

$$\mathbb{F}(\alpha) = \text{card } F(\alpha) \quad \text{for all cardinal numbers } \alpha.$$

This canonical system extends every system of functorial equations solvable by F . If \mathcal{S} is a uniquely solvable system and F is its solution, then the canonical system of F is also uniquely solvable, i.e., F satisfies the following condition:

if $G : \mathbf{SET} \rightarrow \mathbf{SET}$ is a functor with $\text{card } GX = \text{card } FX$ for all sets X , then G is naturally equivalent to F .

The functors satisfying this condition are called DVO-functors (i.e., Determined by their Values on Objects). The DVO-functors are investigated in [2,3,4]. In [4], every DVO-functor is proved to be finitary, which immediately implies that the solution of any uniquely

solvable system of functorial equations is finitary. However, in [4] this result is proved only under a specific set-theoretical hypothesis. The aim of the present paper is to give an absolute (unfortunately, more involved) proof. Here we prove the following (absolute!)

Main Theorem. *Every DVO-functor $\text{SET} \rightarrow \text{SET}$ is finitary.*

Its converse is false, for there are many finitary functors which are not DVO. On the other hand, there are also many finitary functors which are DVO (see [2,3,4]; the full description of all DVO-functors remains unresolved). Hence there also are many uniquely solvable systems of functorial equations: all the canonical systems of the DVO-functors and, possibly, some of their reducts (but a small system of functorial equations, i.e., one consisting only of a set of equations, is never uniquely solvable, see [10]).

Finally, let us mention that the above field of problems can be easily transformed to a more general setting: for arbitrary categories \mathbb{K}, \mathbb{L} a functorial equation

$$F(X) = Y \quad \text{with } X \in \text{obj } \mathbb{K}, \quad Y \in \text{obj } \mathbb{L}$$

is solvable by any functor $F : \mathbb{K} \rightarrow \mathbb{L}$ with FX isomorphic to Y ; the concept of solvability and unique solvability of systems of functorial equations is evident. Also, every functor $F : \mathbb{K} \rightarrow \mathbb{L}$ determines its canonical system of functorial equations; this system is uniquely solvable if and only if F is a DVO-functor (i.e., naturally equivalent to any $G : \mathbb{K} \rightarrow \mathbb{L}$ with GX isomorphic to FX for every $X \in \text{obj } \mathbb{K}$).

Problem. For which cocomplete categories \mathbb{K} and \mathbb{L} is every DVO-functor $\mathbb{K} \rightarrow \mathbb{L}$ finitary?

2. THE IDEA OF THE PROOF AND THE PRELIMINARIES

2.1 The present paper is completely devoted to the proof of Main Theorem. The general scheme of the proof is quite straightforward: given a functor $H : \text{SET} \rightarrow \text{SET}$ which is not finitary, one has to find a functor $G : \text{SET} \rightarrow \text{SET}$, not naturally equivalent to H , such

that $\text{card } GX = \text{card } HX$ for all sets X . In fact, we shall construct two functors $G_1, G_2 : \text{SET} \rightarrow \text{SET}$ which are not naturally equivalent and such that

$$\text{card } HX = \text{card } G_1X = \text{card } G_2X \quad \text{for all sets } X.$$

The reason for doing this is that the internal structure of the given functor H could be very complicated, while only a partial knowledge of it suffices to find many functors $G : \text{SET} \rightarrow \text{SET}$ with $\text{card } HX = \text{card } GX$ for all sets X . But a direct proof that H is not naturally equivalent to such a functor G is a problem. If we construct two such functors G_1, G_2 , both with a relatively simple internal structure, we are able to ensure that they are not naturally equivalent. Then at least one of them is not naturally equivalent to H .

2.2 If H is an endofunctor of a locally finitely presentable category \mathbb{K} , then its finitary part H^f is the left Kan extension of the restriction of H to the category of the finitely presentable objects of \mathbb{K} . Then H^f is really finitary (i.e., it preserves the directed colimits) and it is a subfunctor of H , i.e., there is a ‘canonical’ monotransformation of H^f into H (see e.g. [1]).

Clearly, SET is locally finitely presentable and the finitely presentable objects are just finite sets. Since this paper deals only with endofunctors of SET , we shall use a specific description of the above notions which is more suitable for our computation of the cardinalities.

If $H : \text{SET} \rightarrow \text{SET}$ is a functor, its subfunctor is any functor $G : \text{SET} \rightarrow \text{SET}$ such that $GX \subseteq HX$ for all sets X and Gg is the domain-range restriction of Hg for every mapping $g : X \rightarrow X'$ (thus $Hg(GX) \subseteq GX'$). And the finitary part H^f of H is the subfunctor of H given on a set X by the formula

$$H^f X = \bigcup \{ \text{Im } Hg \mid g : Y \rightarrow X, Y \text{ finite} \}$$

(where $\text{Im } k$ denotes the image of a mapping k in question) and $H^f g$ is just the domain-range restriction of Hg for all mappings $g : X \rightarrow X'$. Since Hg sends the set $H^f X$ into $H^f X'$, this definition is correct.

This set-theoretical description permits us to investigate the sets $HX \setminus H^f X$ and to compute their cardinalities. In fact, the functors G_1 and G_2 mentioned in 2.1, will be constructed (in Section 6 of the present paper) so that H^f is also the finitary part of G_1 and G_2 , and

$$\text{card}(HX \setminus H^f X) = \text{card}(G_1 X \setminus H^f X) = \text{card}(G_2 X \setminus H^f X)$$

for all sets X .

2.3 We have to recall some simple properties of endofunctors of SET.

The trivial functor C_\emptyset (=the constant functor to the empty set) is finitary, hence it does not contradict to Main Theorem and we can restrict ourselves only to non-trivial functors. Any non-trivial endofunctor G of SET sends every non-empty set to a non-empty set and there is a natural transformation

$$\mu : \text{Id} \rightarrow G$$

of the identity functor Id into G . In fact, if $\mathbf{1} = \{*\}$ is a standard one-element set, we choose $a \in G\mathbf{1}$ and for every set X we define $\mu_X : X \rightarrow GX$ by

$$\mu_X(x) = Gv_x(a)$$

where $v_x : \mathbf{1} \rightarrow X$ is the mapping sending $*$ to x .

The transformation μ is either a monotransformation or it factorizes as

$$\text{Id} \rightarrow C_{0,1} \rightarrow G$$

where $C_{0,1}$ is the functor sending \emptyset to \emptyset and all non-empty sets to $\mathbf{1}$.

Every transformation $\tau : C_{0,1} \rightarrow G$ is called a distinguished point of G in [5,8] and $\tau_X(*)$ is a distinguished point of G in GX for every non-empty set X . Clearly, $Gg(\tau_X(*)) = \tau_{X'}(*)$ for every mapping $g : X \rightarrow X'$. Hence every distinguished point $p \in GX$ of G in GX lies in $G^f X$ where G^f denotes the finitary part of G .

If A, B are subsets of a set X and $i_A : A \rightarrow X, i_B : B \rightarrow X$ denote the inclusions, then every $x \in \text{Im } Gi_A \cap \text{Im } Gi_B$ is

a distinguished point of G in GX whenever $A \cap B = \emptyset$ or
 an element of $\text{Im } Gi_{A \cap B}$, where $i_{A \cap B} : A \cap B \rightarrow X$ is the inclusion, whenever $A \cap B \neq \emptyset$ (see [8]).

Hence if $x \in GX$ is not a distinguished point of G in GX (e.g. if $x \in GX \setminus G^f X$), then the system

$$\mathfrak{F}_X^G(x) = \{Z \subseteq X \mid x \in \text{Im } Gi_Z, i_Z : Z \rightarrow X \text{ is the inclusion}\}$$

is a filter on the set X , see [5,8].

2.4 Given a functor $H : \mathbf{SET} \rightarrow \mathbf{SET}$ which is not finitary, the filters just described provide a tool to derive a formula for $\text{card}(HX \setminus H^f X)$ in 3.5. The functors G_1, G_2 mentioned in 2.1-2.2 are constructed in Section 6, and elementary expansions discussed in Section 5 are the building blocks of this construction. Transformation monoids investigated in Section 4 serve to prove that the constructed G_1 and G_2 are not naturally equivalent. Observe that, for any functor $K : \mathbf{SET} \rightarrow \mathbf{SET}$, any set X and any $x \in KX$, the system

$$\mathfrak{M}_X^K(x) = \{g : X \rightarrow X \mid Kg(x) = x\}$$

is a transformation monoid and, if $\nu : K \rightarrow K'$ is a natural equivalence then the transformation monoids $\mathfrak{M}_X^K(x)$ and

$$\mathfrak{M}_X^{K'}(\nu_X(x)) = \{h : X \rightarrow X \mid K'h(\nu_X(x)) = \nu_X(x)\}$$

are strongly isomorphic (for details, see Section 4). Transformation monoids which are not strongly isomorphic are inserted at the appropriate places in the construction of G_1 and G_2 , and this ensures that G_1 and G_2 are not naturally equivalent (for details see Section 6). This will finish our proof.

3. ABSTRACT FILTERS

3.1 Definition. Let \mathcal{F} be a filter on a set X and \mathcal{G} be a filter on a set Y . We say that they are equivalent if there exist $F \in \mathcal{F}, G \in \mathcal{G}$ and a bijection b of F onto G such that, for every $F' \subseteq F$,

$$F' \in \mathcal{F} \quad \text{if and only if} \quad b(F') \in \mathcal{G}.$$

Any class \mathcal{A} of all mutually equivalent filters is called an abstract filter. If a filter \mathcal{F} (on a set X) is an element of an abstract filter \mathcal{A} , we say that \mathcal{F} is a location (on the set X) of the abstract filter \mathcal{A} . Let us denote $\mathcal{A}(X)$ the set of all locations of \mathcal{A} on X .

Remark. By the above equivalence, the class of all filters (on all sets) is decomposed into classes of mutually equivalent filters. Let $|\mathcal{F}|$ denote $\min\{\text{card } F \mid F \in \mathcal{F}\}$. If \mathcal{F} and \mathcal{G} are locations of an abstract filter \mathcal{A} , then, clearly, $|\mathcal{F}| = |\mathcal{G}|$ and $\text{card } \bigcap \mathcal{F} = \text{card } \bigcap \mathcal{G}$. Let us denote $|\mathcal{A}| = |\mathcal{F}|$ and $|\bigcap \mathcal{A}| = \text{card } \bigcap \mathcal{F}$ for a location \mathcal{F} of \mathcal{A} (on a set X).

Observation. If \mathcal{F} is a location of \mathcal{A} on a set X and if $f : X \rightarrow Y$ is a map injective on some $F \in \mathcal{F}$ then the filter \mathcal{G} with a basis $\{f(F) \mid F \in \mathcal{F}\}$ is a location of \mathcal{A} on Y . We shall write $\mathcal{G} = f(\mathcal{F})$.

3.2 Abstract filters and their locations are useful tool for the examination of functors $\text{SET} \rightarrow \text{SET}$ and the following lemma will be often used.

Lemma. *For every abstract filter \mathcal{A} and every set Y , $\mathcal{A}(Y) = \emptyset$ if $\text{card } Y < |\mathcal{A}|$ and $\text{card } \mathcal{A}(Y) \geq \text{card } Y$ if $\text{card } Y \geq \max\{|\mathcal{A}|, \aleph_0\}$.*

Proof. If $F \in \mathcal{F}$ for a location \mathcal{F} of \mathcal{A} then $\text{card } F \geq |\mathcal{A}|$. Hence if \mathcal{F} is a location of \mathcal{A} on a set Y then $\text{card } Y \geq |\mathcal{A}|$. Thus $\mathcal{A}(Y) = \emptyset$ for all sets Y with $\text{card } Y < |\mathcal{A}|$. If $\text{card } Y \geq \max\{|\mathcal{A}|, \aleph_0\}$ then $\text{card}(Y \times Y) = \text{card } Y$ and since on every fibre $Y \times \{y\}$ there is a location of \mathcal{A} , it follows $\text{card } \mathcal{A}(Y) \geq \text{card } Y$. \square

3.3 For a functor H and $x \in HX$ let us recall (see 2.3) the family

$$\mathfrak{F}_X^H(x) = \{Y \subseteq X \mid x \in \text{Im } Hi \text{ for the inclusion } i : Y \rightarrow X\}.$$

If $x \in HX$ is non-distinguished then $\mathfrak{F}_X^H(x)$ is a filter on X .

Clearly, if $x \in HX \setminus H^f X$ then $|\mathfrak{F}_X^H(x)|$ is infinite.

Notation. For an arbitrary functor $H : \mathbf{SET} \rightarrow \mathbf{SET}$ and for a filter \mathcal{F} on a set X , let us denote

$$p(H, \mathcal{F}) = \{x \in HX \mid x \text{ is non-distinguished, } \mathfrak{F}_X^H(x) = \mathcal{F}\}.$$

3.4 Lemma. *Let $H : \mathbf{SET} \rightarrow \mathbf{SET}$ be a functor and let \mathcal{F} and \mathcal{G} be locations of an abstract filter \mathcal{A} on X and Y , respectively. Then there exists a mapping $f : X \rightarrow Y$ such that Hf maps bijectively $p(H, \mathcal{F})$ onto $p(H, \mathcal{G})$. If both \mathcal{F} and \mathcal{G} are locations of \mathcal{A} on a set X and if $\mathcal{F} \neq \mathcal{G}$ then $p(H, \mathcal{F}) \cap p(H, \mathcal{G}) = \emptyset$.*

Proof. If both \mathcal{F} and \mathcal{G} are locations of \mathcal{A} , \mathcal{F} on X and \mathcal{G} on Y , then there exists a bijection b of some $F \in \mathcal{F}$ onto some $G \in \mathcal{G}$ such that for $F' \subseteq F$, $F' \in \mathcal{F}$ if and only if $b(F') \in \mathcal{G}$. If $f : X \rightarrow Y$ is an arbitrary extension of b then $\mathcal{G} = f(\mathcal{F})$ (see 3.1 Observation) and hence $Hf(x) \in p(H, \mathcal{G})$ for all $x \in p(H, \mathcal{F})$, see also [5,9]. Hence Hf maps $p(H, \mathcal{F})$ bijectively onto $p(H, \mathcal{G})$. If $X = Y$ and $x \in p(H, \mathcal{F}) \cap p(H, \mathcal{G})$ then

$$\mathcal{F} = \mathfrak{F}_X^H(x) = \mathcal{G}. \quad \square$$

3.5 Convention. In what follows, the symbol

A

denotes the system of all abstract filters \mathcal{A} with $|\mathcal{A}| \geq \aleph_0$.

By 3.4, we get $\text{card } p(H, \mathcal{F}) = \text{card } p(H, \mathcal{G})$ whenever both \mathcal{F} and \mathcal{G} are locations of an abstract filter \mathcal{A} ; let us denote this cardinal number $p(H, \mathcal{A})$. Then for every $X \neq \emptyset$

$$\text{card}(HX \setminus H^{(f)}X) = \sum_{\mathcal{A} \in \mathbf{A}} p(H, \mathcal{A}) \text{card } \mathcal{A}(X).$$

4. TRANSFORMATION MONOIDS

4.1 Let us recall that a transformation monoid M on a set X is a set of mappings $f : X \rightarrow X$ closed with respect to the composition of mappings and containing the identity mapping. We abbreviate the words ‘transformation monoid’ to ‘ t -monoid’.

If M is a t -monoid on a set X and M' is a t -monoid on a set Y then we say that they are strongly isomorphic if there exists a bijection $b : X \rightarrow Y$ such that

$$f \mapsto b \circ f \circ b^{-1}$$

is a monoid isomorphism of M onto M' .

4.2 For every functor $G : \mathbf{SET} \rightarrow \mathbf{SET}$, every $x \in GX$ determines a t -monoid $\mathfrak{M}_X^G(x)$ on X , namely

$$\mathfrak{M}_X^G(x) = \{f : X \rightarrow X \mid Gf(x) = x\}.$$

If μ is a natural equivalence of G onto a functor G' then, clearly, for every set X and every $x \in GX$,

$$\mathfrak{M}_X^G(x) \text{ is strongly isomorphic to } \mathfrak{M}_X^{G'}(\mu_X(x)).$$

The t -monoids form a more subtle tool for examining set functors than filters (e.g. if $x, y \in GX$ and $\mathfrak{F}_X^G(x) = \mathfrak{F}_X^G(y)$, then not necessarily $\mathfrak{M}_X^G(x) = \mathfrak{M}_X^G(y)$), and we shall use them in our construction.

4.3 For a filter \mathcal{F} on a set X , let $\mathfrak{M}(\mathcal{F})$ denote the t -monoid consisting of $f : X \rightarrow X$ which are injective on a set from \mathcal{F} and $\{f(F) \mid F \in \mathcal{F}\}$ form a basis of \mathcal{F} .

One can verify easily that

- (1) $\mathfrak{M}(\mathcal{F})$ is really a t -monoid on X ;
- (2) if $g \in \mathfrak{M}(\mathcal{F})$ is injective on a set $F \in \mathcal{F}$ and $f : X \rightarrow X$ is a mapping inverse to g on $g(F)$ then $f \in \mathfrak{M}(\mathcal{F})$;
- (3) an idempotent mapping $g : X \rightarrow X$ is in $\mathfrak{M}(\mathcal{F})$ if and only if $\text{Im } g \in \mathcal{F}$;
- (4) $\mathcal{F} = \{\text{Im}(f) \mid f \in \mathfrak{M}(\mathcal{F})\}$.

4.4 Now, let us suppose that $\text{card} \bigcap \mathcal{F} \geq 3$. Let us choose distinct $u, v \in \bigcap \mathcal{F}$ and denote

$$\begin{aligned} \mathfrak{M}(\mathcal{F}, u) &= \{f \in \mathfrak{M}(\mathcal{F}) \mid f(u) = u\} \quad \text{and} \\ \mathfrak{M}(\mathcal{F}, u, v) &= \{f \in \mathfrak{M}(\mathcal{F}) \mid f(u) = u, f(v) = v\}. \end{aligned}$$

Proposition. $\mathfrak{M}(\mathcal{F}, u)$ is not strongly isomorphic to $\mathfrak{M}(\mathcal{F}, u, v)$.

Proof. We prove that $\{x \in X \mid \forall f \in \mathfrak{M}(\mathcal{F}, u), f(x) = x\} = \{u\}$. Since $f(u) = u$ and $f(v) = v$ for all $f \in \mathfrak{M}(\mathcal{F}, u, v)$ the proof will be complete. Consider $x \in X \setminus \bigcap \mathcal{F}$, then $X \setminus \{x\} \in \mathcal{F}$ and therefore every mapping $f : X \rightarrow X$ such that $f(y) = y$ for all $y \in X \setminus \{x\}$ and $f(x) \neq x$ belongs to $\mathfrak{M}(\mathcal{F}, u)$ (and also to $\mathfrak{M}(\mathcal{F}, u, v)$). A mapping f which is an arbitrary permutation of $\bigcap \mathcal{F}$ and $f(y) = y$ for all $y \in X \setminus \bigcap \mathcal{F}$ belongs to $\mathfrak{M}(\mathcal{F})$. Since $\text{card} \bigcap \mathcal{F} \geq 3$ a suitable choice of a permutation guarantees the required statement. \square

Remark. This proposition will be used in the proof of Main Theorem to show that the functors G_1 and G_2 , which we shall construct in 6., are not naturally equivalent.

4.5 In the rest of the paragraph we assume that a filter \mathcal{F} on a set X with $\bigcap \mathcal{F} \neq \emptyset$ is given.

Definition. A mapping $f : X \rightarrow Y$ is called \mathcal{F} -simple if there exists a set $F \in \mathcal{F}$ such that f is injective on F .

Fix a set $\emptyset \neq W \subseteq \bigcap \mathcal{F}$. We write that $f_1 \sim_W f_2$ for \mathcal{F} -simple mappings $f_1, f_2 : X \rightarrow Y$ if there exist $F \in \mathcal{F}$ and $g \in \mathfrak{M}(\mathcal{F})$ such that $g(w) = w$ for all $w \in W$ and $f_1 \circ g(x) = f_2(x)$ for all $x \in F$.

4.6 **Lemma.** For every set Y , the relation \sim_W on the set of all \mathcal{F} -simple mappings $f : X \rightarrow Y$ is an equivalence.

Proof. Clearly, \sim_W is reflexive. We prove that \sim_W is symmetric. Let $f_1, f_2 : X \rightarrow Y$ be \mathcal{F} -simple mappings with $f_1 \sim_W f_2$. Then there exist $g \in \mathfrak{M}(\mathcal{F})$ and $F \in \mathcal{F}$ such that $g(w) = w$ for all $w \in W$ and $f_1 \circ g(x) = f_2(x)$ for all $x \in F$. We can assume that g is injective on F because $F \in \mathcal{F}$ and $g \in \mathfrak{M}(\mathcal{F})$. Then $g(F) \in \mathcal{F}$. By 4.3(2),

there exists $\bar{g} : X \rightarrow X \in \mathfrak{M}(\mathcal{F})$ such that $\bar{g} \circ g(x) = x$ for all $x \in F$. Hence $\bar{g}(w) = w$ for all $w \in W$ because $W \subseteq F$. For every $y \in g(F)$, $f_2 \circ \bar{g}(y) = f_1 \circ g \circ \bar{g}(y) = f_1(y)$ and hence $f_2 \sim_W f_1$. Now we show that \sim_W is transitive. Let $f_1 \sim_W f_2 \sim_W f_3$ for \mathcal{F} -simple mappings $f_1, f_2, f_3 : X \rightarrow Y$. Then there exist $g, g' \in \mathfrak{M}(\mathcal{F})$ and $F, F' \in \mathcal{F}$ such that $g(w) = g'(w) = w$ for all $w \in W$, $f_1 \circ g(x) = f_2(x)$ for all $x \in F$ and $f_2 \circ g'(x) = f_3(x)$ for all $x \in F'$. Then $Z = F' \cap (g')^{-1}(F) \in \mathcal{F}$ and $f_1 \circ (g \circ g')(z) = f_2 \circ g'(z) = f_3(z)$ for all $z \in Z$. Clearly, $g \circ g' \in \mathfrak{M}(\mathcal{F})$ and $g \circ g'(w) = w$ for all $w \in W$. Hence $f_1 \sim_W f_3$. \square

4.7 Lemma. *Let $f_1, f_2 : X \rightarrow Y$ be \mathcal{F} -simple mappings with $f_1 \sim_W f_2$ and let $h : Y \rightarrow Z$ be an arbitrary mapping. Then either both $h \circ f_1$ and $h \circ f_2$ are \mathcal{F} -simple mappings with $h \circ f_1 \sim_W h \circ f_2$ or neither $h \circ f_1$ nor $h \circ f_2$ is \mathcal{F} -simple and $h \circ f_1(w) = h \circ f_2(w)$ for all $w \in W$.*

Proof. We have only to prove that $h \circ f_1$ is \mathcal{F} -simple if and only if $h \circ f_2$ is \mathcal{F} -simple, the other statements are obvious. Since $f_1 \sim_W f_2$ there exist $g \in \mathfrak{M}(\mathcal{F})$ and $F \in \mathcal{F}$ such that $g(w) = w$ for all $w \in W$ and $f_1 \circ g(x) = f_2(x)$ for all $x \in F$. We can assume that g is injective on F . If $h \circ f_1$ is \mathcal{F} -simple then $h \circ f_1$ is injective on a set $F' \in \mathcal{F}$. Consider $F'' = F \cap g^{-1}(F') \in \mathcal{F}$, then $g(F'') \subseteq F'$ and hence $h \circ f_1 \circ g$ is injective on F'' and $h \circ f_1 \circ g(x) = h \circ f_2(x)$ for all $x \in F''$, thus $h \circ f_2$ is \mathcal{F} -simple. By symmetry, we obtain that from the fact that $h \circ f_2$ is \mathcal{F} -simple it follows that $h \circ f_1$ is \mathcal{F} -simple. \square

4.8 Lemma. *Let $f_1 \in \mathfrak{M}(\mathcal{F})$. Then $f_1 \sim_W f_2$ for an \mathcal{F} -simple mapping $f_2 : X \rightarrow X$ if and only if $f_2 \in \mathfrak{M}(\mathcal{F})$ and $f_1(w) = f_2(w)$ for all $w \in W$.*

Proof. Observe that if a mapping $f_2 : X \rightarrow X$ is \mathcal{F} -simple and $f_2 \sim_W f_1$ for $f_1 \in \mathfrak{M}(\mathcal{F})$ then $f_2 \in \mathfrak{M}(\mathcal{F})$ (because $\mathfrak{M}(\mathcal{F})$ is closed under composition) and $f_2(w) = f_1(w)$ for all $w \in W$. Conversely, assume that $f_1, f_2 \in \mathfrak{M}(\mathcal{F})$ such that $f_1(w) = f_2(w)$ for all $w \in W$. Then there exist $F_1, F_2 \in \mathcal{F}$ such that f_i is injective on F_i and $f_i(F_i) \in \mathcal{F}$ for $i = 1, 2$. Then $F = f_1(F_1) \cap f_2(F_2) \in \mathcal{F}$ and also $F'_i = F_i \cap f_i^{-1}(F) \in \mathcal{F}$ for $i = 1, 2$. By 4.3(2), there exists $g' \in \mathfrak{M}(\mathcal{F})$

such that $f_1 \circ g'(x) = x$ for all $x \in F$. Clearly $g = g' \circ f_2 \in \mathfrak{M}(\mathcal{F})$ and $f_1 \circ g(x) = f_1 \circ g' \circ f_2(x) = f_2(x)$ for all $x \in F'_2$. Since $W \subseteq \bigcap \mathcal{F} = g'(\bigcap \mathcal{F}) \subseteq F'_1 \cap F'_2$ and since f_1 and f_2 are one-to-one on $\bigcap \mathcal{F}$ and $f_1(w) = f_2(w)$ for all $w \in W$ we conclude that $g(w) = w$ for all $w \in W$. Thus $f_1 \sim_W f_2$ and the proof is complete. \square

As a consequence we obtain this

Corollary. *The cardinal number of the set $\mathfrak{M}(\mathcal{F}) / \sim_W$ is equal to the cardinal number of the set of all injective mappings from W into $\bigcap \mathcal{F}$.*

5. EXPANSION OF FUNCTORS

5.1 Let $K : \text{SET} \rightarrow \text{SET}$ be a functor and X be a set with $\text{card } X > 1$. We are going to construct a functor G which extends K by the addition of one element, say a , to KX . The functor G has to enclose K and a together ‘as tightly as possible’, i.e., to add new elements to any KY only when it is absolutely necessary, for, in a ‘tight enough’ extension, we shall be able to control the cardinalities of GY . Moreover, we also need to control the internal structure of G , i.e., the knowledge of the filters and of the t -monoids of the newly added elements. This will be possible whenever the filter and the t -monoid of a in GX are prescribed. However, the filter and the t -monoid will have to have properties which make the whole construction possible.

5.2 So let a filter \mathcal{F} on the set X be given such that $|\mathcal{F}| = \text{card } X$, $\bigcap \mathcal{F} \neq \emptyset$. Moreover, let a non-empty set $W \subseteq \bigcap \mathcal{F}$ be given. Recall the t -monoid $\mathfrak{M}(\mathcal{F})$ defined by \mathcal{F} in 4.3, \mathcal{F} -simple mappings $f : X \rightarrow Y$ and the equivalence \sim_W both defined in 4.5. We need them in our construction. We ‘add $Gf(a)$ to KY ’ for every \mathcal{F} -simple mapping $f : X \rightarrow Y$. On the other hand, we want to map $Gf(a)$ into KY whenever $f : X \rightarrow Y$ is not \mathcal{F} -simple. To do it ‘functorially’, we need further instruments: a natural transformation $\mu : \text{Id} \rightarrow K$ of the identity functor Id into K (such μ does exist, see 2.3) and an element $u \in W$. Hence our construction will depend on the quadruple

of ‘parameters’

$$(\mu, \mathcal{F}, W, u).$$

5.3 Construction. For an \mathcal{F} -simple mapping $f : X \rightarrow Y$, let $[f]$ denote the equivalence class of \sim_W on the set of all \mathcal{F} -simple mappings $X \rightarrow Y$ containing f .

For a set Y , define

$$GY = KY \cup \{[f] \mid f : X \rightarrow Y \text{ is } \mathcal{F}\text{-simple}\},$$

where we suppose that the union is disjoint. If $h : Y \rightarrow Z$ is a mapping then for every $y \in GY$ define

$$Gh(y) = \begin{cases} Kh(y) & \text{if } y \in KY, \\ [h \circ f] & \text{if } y = [f] \text{ for } \mathcal{F}\text{-simple } f : X \rightarrow Y \\ & \text{and } h \circ f \text{ is } \mathcal{F}\text{-simple,} \\ Kh(\mu_Y(f(u))) & \text{if } y = [f] \text{ for } \mathcal{F}\text{-simple } f : X \rightarrow Y \\ & \text{and } h \circ f \text{ is not } \mathcal{F}\text{-simple.} \end{cases}$$

Observation. By 4.6 and 4.7, G is a correctly defined functor from \mathbf{SET} into itself and K is its subfunctor and the element a mentioned in 5.1 is precisely $[1_X]$, where 1_X is the identity mapping of X . We call it the elementary expansion of K (determined by (μ, \mathcal{F}, W, u)).

5.4 In the lemmas below $K, X, \mathcal{F}, W, u, \mu$ are as above. Moreover, let \mathcal{A} denote the abstract filter of \mathcal{F} (i.e., \mathcal{F} is a location of \mathcal{A} on the set X , see 3.1).

Lemma. $\mathfrak{F}_Y^G(y)$ is a location of \mathcal{A} for every $y \in GY \setminus KY$ and for every set Y . Further, $\mathfrak{F}_X^G([f]) = \mathcal{F}$ if and only if $f \in \mathfrak{M}(\mathcal{F})$. Moreover, $\mathfrak{M}_X^G([1_X]) = \{f \in \mathfrak{M}(\mathcal{F}) \mid f(w) = w \text{ for all } w \in W\}$.

Proof. Assume that $y = [f]$ for an \mathcal{F} -simple mapping $f : X \rightarrow Y$. Thus there exists $F \in \mathcal{F}$ such that f is injective on F . Consider a set $Z \in f(\mathcal{F})$ then $F' = F \cap f^{-1}(Z) \in \mathcal{F}$. Let $\iota : Z \rightarrow Y$ be

the inclusion mapping, then there exists a mapping $g : X \rightarrow Z$ such that $f(z) = \iota \circ g(z)$ for all $z \in F'$. Since f is \mathcal{F} -simple we conclude that g is \mathcal{F} -simple. By 4.3(3), every idempotent mapping $h : X \rightarrow X$ with $\text{Im}(h) = F'$ belongs to $\mathfrak{M}(\mathcal{F})$, hence $f \sim_W \iota \circ g$ and thus $f(\mathcal{F}) \subseteq \mathfrak{F}_Y^G([f])$. Conversely, if $Z \in \mathfrak{F}_Y^G([f])$ and if $\iota : Z \rightarrow Y$ is the inclusion then there exists an \mathcal{F} -simple mapping $g : X \rightarrow Z$ such that $\iota \circ g \sim_W f$ and hence there exists $F' \in \mathcal{F}$ with $F' \subseteq F$ and $f(F') \subseteq Z$. Therefore $f(\mathcal{F}) = \mathfrak{F}_Y^G([f])$. The fact that f is \mathcal{F} -simple demonstrates that $\mathfrak{F}_Y^G([f])$ is a location of \mathcal{A} . From the definition of $\mathfrak{M}(\mathcal{F})$ it follows that $f(\mathcal{F}) = \mathcal{F}$ for a \mathcal{F} -simple mapping if and only if $f \in \mathfrak{M}(\mathcal{F})$, and the second statement follows. The third statement is implied by Lemma 4.8. \square

5.5 Lemma. *If $W = \{u\}$ and $\text{card} \bigcap \mathcal{F} \geq 3$ then for every set Y and every $y \in GY \setminus KY$, the set $\{z \in Y \mid f(z) = z \text{ for all } f \in \mathfrak{M}_Y^G(y)\}$ is a singleton.*

Proof. Consider $y = [g] \in GY \setminus KY$ and let $U = \{z \in Y \mid f(z) = z \text{ for all } f \in \mathfrak{M}_Y^G([g])\}$. If $h \in \mathfrak{M}_Y^G([g])$ then $Gh([g]) = [g]$ implies that $h \circ g \sim_W g$ and from the definition of \sim_W it follows that $h(g(u)) = g(u)$. Therefore $g(u) \in U$. By Lemma in 5.4, $\mathfrak{F}_Y^G([g])$ is a location of \mathcal{A} and hence $\text{card} \bigcap \mathfrak{F}_Y^G(y) = \text{card} \bigcap \mathcal{F} \geq 3$. One can easily see that if $t \in Y \setminus \bigcap \mathfrak{F}_Y^G([g])$ then the mapping $h : Y \rightarrow Y$ such that $h(t) \neq t$ and $h(s) = s$ for all $s \in Y$ with $s \neq t$ satisfies $Gh([g]) = [g]$ and hence $t \notin U$ (see also [5,9]). If $t \in \bigcap \mathfrak{F}_Y^G([g])$ with $t \neq g(u)$, then there exists $t' \in \bigcap \mathcal{F}$ with $g(t') = t$ and $t' \neq u$. Let $h : X \rightarrow X$ be a mapping such that $h(x) = x$ for all $x \in X \setminus \bigcap \mathcal{F}$, the restriction of h on $\bigcap \mathcal{F}$ is a permutation of $\bigcap \mathcal{F}$ with $h(u) = u$ and $h(t') \neq t'$. Since g is \mathcal{F} -simple there exists $F \in \mathcal{F}$ such that g is injective on F and therefore there exists a mapping $h' : Y \rightarrow Y$ such that $g \circ h(x) = h' \circ g(x)$ for all $x \in F$. Hence $h'(t) \neq t$ and $g \sim_W h' \circ g$. Thus $Gh'([g]) = [g]$ and $t \notin U$. \square

5.6 Summary. *Let $K : \text{SET} \rightarrow \text{SET}$ be a functor, let G be an elementary expansion of K determined by the quadruple (μ, \mathcal{F}, W, u) , and let \mathcal{A} be the abstract filter containing \mathcal{F} . Then*

(1) if $|\mathcal{F}|$ is infinite and W is finite then

$$\text{card}(GY \setminus KY) = \text{card } \mathcal{A}(Y)$$

for every set Y whenever $\mathfrak{F}_Y^K(y)$ is a location of \mathcal{A} for no $y \in KY$;

- (2) there exists $a \in GX \setminus KX$ such that $\mathfrak{M}_X^G(a) = \{f \in \mathfrak{M}(\mathcal{F}) \mid f(w) = w \text{ for all } w \in W\}$;
- (3) if $\text{card } \bigcap \mathcal{F} \geq 3$ and $W = \{u\}$ then for every set Y and every $y \in GY \setminus KY$, $\mathfrak{M}_Y^G(y)$ has exactly one fix-point (i.e., there exists exactly one $v \in Y$ with $f(v) = v$ for all $f \in \mathfrak{M}_Y^G(y)$).

Proof. If $\mathfrak{F}_Y^K(y)$ is a location of \mathcal{A} for no $y \in KY$ then, by 3.5 and 5.4, $\text{card}(GY \setminus KY) = p(G, \mathcal{A}) \text{card } \mathcal{A}(Y)$. By Lemma and Corollary in 4.8,

$$p(G, \mathcal{A}) = \text{card}(\bigcap \mathcal{F})^W \leq |\mathcal{A}|$$

because $\mathcal{A} \in \mathbb{A}$. From 3.2 and $|\mathcal{A}| \geq \aleph_0$ it follows that

$$p(G, \mathcal{A}) \text{card } \mathcal{A}(Y) = \text{card } \mathcal{A}(Y)$$

and (1) is proved. Lemma 5.4 implies (2) and Lemma 5.5 implies (3). \square

6. THE CONSTRUCTION OF G_1 AND G_2

6.1 An amalgam $\mathfrak{A} = \{G^{(j)} \mid j \in J\}$ of functors with a base K is a system of functors such that K is a subfunctor of G_j for all $j \in J$ and

$$G^{(j_1)}X \cap G^{(j_2)}X = KX \text{ for all sets } X \text{ and all } j_1, j_2 \in J \text{ with } j_1 \neq j_2.$$

If, for every set X , $\bigcup_{j \in J} G^{(j)}X$ is a set, we can define the sum of the amalgam \mathfrak{A} by the simple rule

$$GX = \bigcup_{j \in J} G^{(j)}X \text{ and each } G^{(j)} \text{ is a subfunctor of } G.$$

Clearly, G is a correctly defined functor and, for every set X ,

$$\text{card}(GX \setminus KX) = \sum_{j \in J} \text{card}(G^{(j)}X \setminus KX).$$

6.2 Now we are going to complete the proof of Main Theorem. Let a functor $H : \mathbb{SET} \rightarrow \mathbb{SET}$ which is not finitary be given. Then, by 3.5,

$$\text{card}(HX \setminus H^{(f)}X) = \sum_{\mathcal{A} \in \mathbb{A}} p(H, \mathcal{A}) \text{card } \mathcal{A}(X),$$

where \mathcal{A} and $\mathcal{A}(X)$ are as in 3.1, $p(H, \mathcal{A})$ and \mathbb{A} are as in 3.5. Since H is not finitary, $p(H, \mathcal{A}) \neq 0$ for at least one $\mathcal{A} \in \mathbb{A}$.

We aim to construct functors $G_1, G_2 : \mathbb{SET} \rightarrow \mathbb{SET}$ which are not naturally equivalent and satisfy

$$G_1^f = H^f = G_2^f \quad \text{and} \\ \text{card}(G_1X \setminus G_1^fX) = \sum_{\mathcal{A} \in \mathbb{A}} p(H, \mathcal{A}) \text{card } \mathcal{A}(X) = \text{card}(G_2X \setminus G_2^fX)$$

for all sets X . Both G_1 and G_2 will be obtained as sums of suitable amalgams with a base H^f . These amalgams consist of suitable elementary expansions $G_1^{(j)}$ and $G_2^{(j)}$ of H^f . However, to get the quadruples (μ, \mathcal{F}, W, u) from which the elementary expansions will be constructed (see Section 5), we need one more simple trick. For any filter \mathcal{F} on a set X with $|\mathcal{F}| \geq \aleph_0$, put

$$\Phi\mathcal{F} = \begin{cases} \mathcal{F} & \text{if } \bigcap \mathcal{F} \text{ is infinite,} \\ \{F \cup Q \mid F \in \mathcal{F}\} & \text{if } \bigcap \mathcal{F} \text{ is finite (including } \bigcap \mathcal{F} = \emptyset), \end{cases}$$

where Q is a set with $\text{card } Q = 3$ and $X \cap Q = \emptyset$. Clearly, if \mathcal{F} is equivalent (in the sense of 3.1) to \mathcal{G} then $\Phi\mathcal{F}$ is equivalent to $\Phi\mathcal{G}$; hence we have determined $\Phi\mathcal{A}$ for every abstract filter \mathcal{A} and $\text{card } \bigcap \mathcal{F} \geq 3$ for every location \mathcal{F} of $\Phi\mathcal{A}$.

Lemma. *If $\mathcal{A} \in \mathbb{A}$ then*

$$\text{card } \mathcal{A}(Y) = \text{card } \Phi\mathcal{A}(Y) \quad \text{for all sets } Y.$$

Proof. Since $\mathcal{A} \in \mathbb{A}$, $|\mathcal{A}|$ is infinite. If $\text{card } Y < |\mathcal{A}|$ then $\text{card } \mathcal{A}(Y) = 0 = \text{card } \Phi\mathcal{A}(Y)$. If $\text{card } Y \geq |\mathcal{A}| = |\Phi\mathcal{A}|$ then, clearly,

$$\text{card } \Phi\mathcal{A}(Y) \leq \text{card } \mathcal{A}(Y) \text{ card } Y^3.$$

Since $\text{card } Y^3 = \text{card } Y \leq \text{card } \mathcal{A}(Y)$, see 3.2, we conclude that

$$\text{card } \Phi\mathcal{A}(Y) \leq \text{card } \mathcal{A}(Y).$$

The reverse inequality is evident. \square

6.3 Now we are ready to describe the quadruples used in Section 5. First we choose a natural transformation μ from the identity functor to $H^{(f)}$. For every $\mathcal{A} \in \mathbb{A}$, choose one location \mathcal{F} of $\Phi\mathcal{A}$ on a set X with $\text{card } X = |\mathcal{A}|$ and two distinct elements $u, v \in \bigcap \mathcal{F}$. Let $G_1^{\mathcal{A}}$ be the elementary expansion of $H^{(f)}$ determined by the quadruple $(\mu, \mathcal{F}, \{u\}, u)$ and $G_2^{\mathcal{A}}$ be the elementary expansion of $H^{(f)}$ determined by the quadruple $(\mu, \mathcal{F}, \{u, v\}, u)$. Let us denote $p(H, \mathcal{A}) \cdot G_i^{\mathcal{A}}$ the sum of the amalgam of $\mathfrak{A}_i = \{G_i^{(j)} \mid j \in J\}$ for $i = 1, 2$ where $\text{card } J = p(H, \mathcal{A})$, $G_i^{(j)}$ is naturally equivalent to the elementary expansion $G_i^{\mathcal{A}}$ of $H^{(f)}$ for all $j \in J$ and $i = 1, 2$ and $G_i^{(j)} X \cap G_i^{(j')} X = H^{(f)} X$ for all distinct $j, j' \in J$, for all sets X and for $i = 1, 2$. Then, by 5.6, for every set Y and $i = 1, 2$,

$$\begin{aligned} \text{card}((p(H, \mathcal{A}) \cdot G_i^{\mathcal{A}})Y \setminus H^{(f)}Y) &= p(H, \mathcal{A}) \text{card } \Phi\mathcal{A}(Y) = \\ &= p(H, \mathcal{A}) \text{card } \mathcal{A}(Y). \end{aligned}$$

Finally, let G_i be the sum of the amalgam $\{p(H, \mathcal{A}) \cdot G_i^{\mathcal{A}} \mid \mathcal{A} \in \mathbb{A}\}$, for $i = 1, 2$. Then, for every set Y and for $i = 1, 2$,

$$\text{card } G_i Y = \text{card } H^{(f)} Y + \sum_{\mathcal{A} \in \mathbb{A}} p(H, \mathcal{A}) \text{card } \mathcal{A}(Y) = \text{card } H Y,$$

by the equation in 3.5.

6.4 It remains to show that G_1 is not naturally equivalent to G_2 . Since $H \neq H^{(f)}$, there exists $\mathcal{A}_0 \in \mathbb{A}$ such that $p(H, \mathcal{A}_0) \neq 0$. Let \mathcal{F} be a location of $\Phi\mathcal{A}_0$ on a set X with $\text{card } X = |\mathcal{A}_0|$. Assume that ν is a natural equivalence of G_1 onto G_2 . Then ν maps the finitary part $H^{(f)}$ of G_1 onto the finitary part $H^{(f)}$ of G_2 , hence ν_X maps $G_1X \setminus H^{(f)}X$ bijectively onto $G_2X \setminus H^{(f)}X$. Then for every $x \in G_1X \setminus H^{(f)}X$, the t -monoid $\mathfrak{M}_X^{G_1}(x)$ must be strongly isomorphic to $\mathfrak{M}_X^{G_2}(\nu_X(x))$, see 4.2. But for every $x \in G_1X$, the t -monoid $\mathfrak{M}_X^{G_1}(x)$ has at most one fix-point, see 5.5, and $\mathfrak{M}_X^{G_2}[1_X]$ has at least two fix-points, u and v . This is a contradiction, and therefore G_1 and G_2 are not naturally equivalent.

The proof of Main Theorem is now complete.

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