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Universal Codes and Unimodular Lattices

par Robin CHAPMAN et Patrick SOLÉ

Résumé. Les codes résidus quadratiques binaires de longueur $p+1$ produisent par construction B et bourrage des réseaux de type II comme le réseau de Leech. Récemment, il a été prouvé que les codes résidus quadratiques quaternaires produisent les mêmes réseaux par construction A modulo 4. Nous montrons de manière directe l' équivalence des deux constructions pour $p \leq 31$. En dimension 32 nous obtenons un réseau extrémal de type II qui n'est pas isomètre au réseau de Barnes-Wall BW_{32}. On considère également l'équivalence entre construction B modulo 4 plus bourrage et construction A modulo 8. En dimension 48 elles conduisent toutes deux à une nouvelle description du réseau extrémal de type II appelé P_{48q}.

Abstract. Binary quadratic residue codes of length $p+1$ produce via construction B and density doubling type II lattices like the Leech. Recently, quaternary quadratic residue codes have been shown to produce the same lattices by construction A modulo 4. We prove in a direct way the equivalence of these two constructions for $p \leq 31$. In dimension 32, we obtain an extremal lattice of type II not isometric to the Barnes-Wall lattice BW_{32}. The equivalence between construction B modulo 4 plus density doubling and construction A modulo 8 is also considered. In dimension 48 they both led to a new description of the extremal type II lattice P_{48q}.

1. Introduction

In [2], Bonnecaze, Solé and Calderbank introduce for primes $p \equiv \pm 1 \pmod 8$, codes $\widehat{\mathcal{Q}}$ and $\widehat{\mathcal{N}}$, the *universal extended quadratic residue codes*, of length $p+1$ over the 2-adic integers Z_{2^∞}. For positive integers s they consider their reductions $\widehat{\mathcal{Q}}_{2^s}$ and $\widehat{\mathcal{N}}_{2^s}$ modulo 2^s; $\widehat{\mathcal{Q}}_2$ and $\widehat{\mathcal{N}}_2$ are just the standard binary extended quadratic residue codes, while $\widehat{\mathcal{Q}}_4$ and $\widehat{\mathcal{N}}_4$ are the *quaternary quadratic residue codes*. Given a code C of length n over Z_4

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define $\Lambda(C)$ as the set of vectors in Z^n which reduce modulo 4 to elements of C. If $p \equiv -1 \pmod 8$ the lattice $\frac{1}{2}\Lambda(\widehat{Q}_4)$ is even and unimodular ([2] Corollary 4.1); if $p = 7$ it is the E_8 lattice, while if $p = 23$ it is the Leech lattice.

Here we show, by means of an explicit isomorphism, that if $p \equiv -1 \pmod 8$ and $p \leq 31$ then $\frac{1}{2}\Lambda(\widehat{Q}_4)$ is isometric to a lattice $L(\widehat{Q}_2)$ constructed from the binary quadratic residue code in a manner (construction B plus density doubling) generalizing the original construction of the Leech lattice. If $p = 23$ this yields a short proof of what is perhaps the simplest construction of the Leech lattice [2]. If $p = 31$ this, combined with results of Koch and Venkov, shows that $BSBM_{32}$ introduced in [1] is not isometric to the Barnes-Wall lattice BW_{32} . In section §4 we consider a quaternary analogue of this situation, replacing construction B by construction B modulo 4, and construction A mod 4 by construction A mod 8. We show, inter alia, that P_{48_q} can be obtained in the latter way from a quadratic residue code of length 48 over Z_8.

2. The main result

Throughout this section we assume that p is a prime satisfying $p \equiv -1 \pmod 8$. We also fix an integer r such that $r \equiv 1 \pmod 4$, and $r^2 + p \equiv 0 \pmod{32}$. In addition if $p \leq 31$ we will assume that $r^2 + p = 32$. (If $p = 7$, 23 or 31, then $r = 5$, -3 or 1 respectively.)

We first outline a construction of lattices from binary codes of length $p+1$. Consider a self-orthogonal linear subcode C of Z_2^{p+1}, containing the all-ones word. Define $L(C)$ to be the sublattice of Z^{p+1} generated by the following types of vectors:

(1) all vectors of shape $(8\ 0^p)$,
(2) all vectors of shape $(4^2\ 0^{p-1})$,
(3) all vectors of shape $(2^a\ 0^{p+1-a})$ whose support coincides with the support of an element of C,
(4) any vector of shape $(r\ 1^p)$.

This can be recast as the union of two cosets

$$2B(C) \cup ((r\ 1^p) + 2B(C)),$$

of the lattice $2B(C)$ obtained, up to scaling, by construction B applied to C namely

$$B(C) := C + 2P_{p+1} + 4Z^{p+1},$$

with P_{p+1} denoting the parity-check code of length $p+1$. It is clear that $L(C)$ is a lattice, of index $4^{p+1}/|C|$ in Z^{p+1}. If the code C is doubly even, then the norm of each vector in $L(C)$ is divisible by 16. It follows that if C is self-dual and doubly even then the lattice $\frac{1}{\sqrt{8}}L(C)$ is even and unimodular. If C is $\widehat{\mathcal{Q}}_2$ or $\widehat{\mathcal{N}}_2$ then it has these properties. We give four examples of this construction.

p	lattice	reference
7	Gosset	[7]
23	Leech	[6, p.131]
31	BW_{32}	[7, 9]
31	$BSBM_{32}$	[1]

By the *norm* of an element in Euclidean space we mean the square of its length, and the *minimal norm* of a lattice is the least norm of a non-zero element of the lattice. It is easy to see that the minimum norm of $\frac{1}{\sqrt{8}}L(\widehat{\mathcal{Q}}_2)$ is $\min(4, 2\lceil\frac{p+1}{16}\rceil, \frac{1}{2}\mathrm{mw}(\widehat{\mathcal{Q}}_2))$ where $\mathrm{mw}(C)$ is the minimum (Hamming) weight of the code C.

THEOREM 1. *The lattices* $\frac{1}{2}\Lambda(\widehat{\mathcal{Q}}_4)$ *and* $\frac{1}{\sqrt{8}}L(\widehat{\mathcal{Q}}_2)$ *are isometric for* $p \leq 31$.

Proof. Assume $p \leq 31$. We recall the definition of $\widehat{\mathcal{Q}}$ from §III of [2]. Let δ be the square root of $-p$ in Z_{2^∞} with $\delta \equiv -1 \pmod 4$. Note then that $\delta \equiv -r \pmod{16}$. The vectors m_α ($\alpha \in F_p \cup \{\infty\}$) are defined as the rows of the matrix

$$M = \begin{pmatrix} \delta & 1 \cdots 1 \\ -1 & \\ \vdots & W + \delta I \\ -1 & \end{pmatrix}$$

where

$$W_{ij} = \left(\frac{j-i}{p}\right).$$

(The rows and columns of this matrix are labelled in the order $\infty, 0, 1, \dots, p-1$.) The matrix W is called a Jacobsthal matrix , and is instrumental in building Hadamard matrices of Paley type [10, Chap. II]. We collect here the properties that we need

(J1) $JW = WJ = 0$
(J2) $WW^T = pI - J$
(J3) $A := \sum_{i=\square} W_{-i,1} = -1$
(J4) $B := \sum_{i=\square} W_{i,1} = 0$

where J stands for the all-one matrix. See [10, Chap. II, Lemma 7] for proofs of (J1) and (J2). To prove (J3), (J4) observe firstly that by ($J1$) we have, knowing that -1 is not a quadratic residue, that $A + B = -1$.

Secondly, writing χ for the Jacobi symbol we have

$$B = \frac{1}{2} \sum_{x \in F_p,\, x \neq 0} \chi(1 - x^2)$$

and by the character property of χ

$$B = \frac{1}{2} \sum_{x \in F_p,\, x \neq 0} \chi(1 - x)\chi(1 + x) = 0,$$

the last equality coming from (J2).

The coordinate positions in the code are labelled $\infty, 0, 1, \dots, p-1$, regarded as elements of the projective line over F_p. The universal extended quadratic residue code is now defined as

$$\widehat{Q} = \left(\sum_{\alpha \in F_p \cup \{\infty\}} Q_{2^\infty} m_\alpha \right) \cap Z_{2^\infty}^{p+1},$$

where Q_{2^∞} is the field of 2-adic numbers. A similar definition holds for $\widehat{\mathcal{N}}$ with $W_{i,j}$ replaced by $W_{i,-j}$.

We can now describe $\Lambda(\widehat{Q}_4)$ as the set of vectors in Z^{p+1} congruent modulo 4 to elements of $\widehat{Q}$. Let $n_\alpha \in Z^{p+1}$ be the rows of the matrix

$$N = \begin{pmatrix} -r & 1 \cdots 1 \\ -1 & \\ \vdots & W - rI \\ -1 & \end{pmatrix}$$

so that for $\alpha \in F_p \cup \{\infty\}$ we have $n_\alpha \equiv m_\alpha \pmod{16}$. Since by (J2) we have $NN^t = 32I$ the matrix $\frac{1}{\sqrt{32}}N$ is orthogonal. We claim that this matrix maps $\frac{1}{\sqrt{8}}L(\widehat{\mathcal{N}_2})$ to $\frac{1}{2}\Lambda(\widehat{Q}_4)$. This is equivalent to saying that N maps $\frac{1}{8}L(\widehat{\mathcal{N}_2})$ to $\Lambda(\widehat{Q}_4)$. Note that these codes and lattices are preserved by the automorphism σ coming from the permutation $(0\ 1\ 2\ \cdots\ p-1)$ on $F_p \cup \{\infty\}$, and this automorphism maps m_α to $m_{\alpha+1}$ and n_α to $n_{\alpha+1}$. This automorphism is the shift in the cyclic construction of the QR codes. We proceed to show that the images by N of the four types of vectors in construction L above lie in $\Lambda(\widehat{Q}_4)$.

Since the $n_\alpha \in \Lambda(\widehat{Q}_4)$, the matrix N takes the coordinate vectors, which lie in $\frac{1}{8}L(\widehat{\mathcal{N}_2})$, into $\Lambda(\widehat{Q}_4)$. For convenience let $(a, b; c; d)$ denote the vector with ∞-coordinate a, 0-coordinate b, and generic α-coordinate c, and generic β-coordinate d where α and β are any quadratic residue, and quadratic non-residue respectively. Now

$$\frac{1}{2}(n_\infty + n_0) = \left(\frac{r-1}{2}, \frac{-r-1}{2}; 0; -1\right)$$

which lies in Z^{p+1} and is congruent to $\frac{1}{2}(m_0 - m_\infty)$ modulo 8. Hence $\frac{1}{2}(m_\infty + m_0) \in \Lambda(\widehat{Q}_4)$. Applying σ it follows that $\frac{1}{2}(m_\infty + m_\alpha) \in \Lambda(\widehat{Q}_4)$ for all $\alpha \in F_p$, and so $\frac{1}{2}(m_\alpha + m_\beta) \in \Lambda(\widehat{Q}_4)$ for all $\alpha, \beta \in F_p \cup \{\infty\}$. Hence $\frac{1}{8}vN \in \Lambda(\widehat{Q}_4)$ for all v of the shape $(4^2\ 0^{p-1})$. We next compute

$$\frac{1}{4}\left(n_\infty + \sum_{j\in Q'} n_j\right) = \left(-\frac{2r+p-1}{8}, -\frac{p+1}{8}; 0; -\frac{r-1}{4}\right)$$

where Q' is the set of quadratic non-residues modulo p. The last coordinates estimates come from (J3), (J4). Again this has integer coordinates, and is congruent modulo 4 to $\frac{1}{4}(m_\infty + \sum_{j=\square} m_j)$, so this vector lies in $\Lambda(\widehat{Q}_4)$. It follows that

$$\frac{1}{4}\left(n_\infty + \sum_{j\in Q'} n_{j+k}\right) \in \Lambda(\widehat{Q}_4)$$

for each $k \in F_p$. But $\widehat{\mathcal{N}_2}$ is generated by the vectors whose supports are the sets $\{\infty\} \cup (k + Q')$. ([2, p.370, III. A.]). It follows that if v has the shape $(2^a\ 0^{p+1-a})$ and whose support is the same as that of an element of $\widehat{\mathcal{N}_2}$, then $\frac{1}{8}vN \in \Lambda(\widehat{Q}_4)$. Finally

$$\frac{1}{8}\left(rn_\infty + \sum_{j\in F_p} n_j\right) = (-4, 0; 0; 0) \in \Lambda(\widehat{Q}_4)$$

and so $\frac{1}{8}(r, 1; 1; 1)N \in \Lambda(\widehat{Q}_4)$. Hence $\frac{1}{8}L(\widehat{\mathcal{N}_2})N \subseteq \Lambda(\widehat{Q}_4)$, and comparing determinants we see that $\frac{1}{8}L(\widehat{\mathcal{N}_2})N = \Lambda(\widehat{Q}_4)$. Since $L(\widehat{Q}_2)$ and $L(\widehat{\mathcal{N}_2})$ are isometric the Theorem follows. □

3. Application to the cases of $p = 23, 31$.

If $(a_1, \dots, a_n)$ is an element of a code over Z_4, then its *Euclidean weight* is $w(a_1) + \cdots + w(a_n)$ where

$$w(a) = \begin{cases} 0 & \text{if } a = 0, \\ 1 & \text{if } a = \pm 1, \\ 4 & \text{if } a = 2. \end{cases}$$

The *minimum Euclidean weight* mew(C) of a code C over Z_4 is the least Euclidean weight of its non-zero elements. If C is a linear code then the minimum norm of $\Lambda(C)$ is $\min(16, \mathrm{mew}(C))$. For $p = 23$ and $p = 31$, the minimum norm of $L(\widehat{\mathcal{N}_2})$ is 32, and so the minimum norm of $\Lambda(\widehat{\mathcal{Q}}_4)$ is 16. Hence $\mathrm{mew}(\widehat{\mathcal{Q}}_4) \geq 16$. In [2] this is proved in a more elaborate way for $p = 23$.

In [8] Koch and Venkov show that for the five non-isomorphic doubly even self-dual binary codes $C_1, \ldots, C_5$ of length 32, the lattices $L(C_1), \ldots, L(C_5)$ are all non-isometric. We can take $C_1 = \widehat{\mathcal{Q}}_2$, and C_2 to be the Reed-Muller code $RM(2,5)$. Since $L(RM(2,5))$ is isometric to the Barnes-Wall lattice BW_{32} [9], it follows that $\frac{1}{2}\Lambda(\widehat{\mathcal{Q}}_4)$ for $p = 31$ is not isometric to BW_{32}, confirming a conjecture of [1]. It is known that there are only two unimodular lattices in dimension 32 with minimal norm 4 and an automorphism of order 31 [12]. From the results of [1] and of the current paper we can infer than both can be constructed by construction A mod 4 applied to an extended quaternary cyclic code: the quaternary Reed-Muller code $QRM(2,5)$ in the case of BW_{32} and the extended quadratic residue code $\widehat{\mathcal{Q}}_4$ in the case of $BSBM_{32} := \frac{1}{2}\Lambda(\widehat{\mathcal{Q}}_4)$. Both lattices also appear in [11, 4].

4. Quaternary Analogue

We assume in this § that $p \geq 47$ is a prime $\equiv -1 \pmod 8$, and that the integer $r \equiv 1 \pmod 4$ satisfies

$$r^2 + p = 96 = 16.6,$$

if $p = 47, 71$ and

$$r^2 + p = 128 = 16.8.$$

if $p = 79, 103, 127$. The corresponding values of r are $r = -7, 5$ in first case and $r = -7, 5, 1$ in the second. For a quaternary code C of length $p+1$ we define

$$B_4(C) := C + 4P_{p+1} + 8Z^{p+1},$$

and

$$L_4(C) := 2B_4(C) \cup ((r\ 1^p) + 2B_4(C)).$$

For an octonary code C_8 of length $p+1$, we define

$$\Lambda_4(C_8) = C_8 + 8Z^{p+1}.$$

We have the following analogue of Theorem 1:

THEOREM 2. *The lattices $\frac{1}{4}L_4(\widehat{\mathcal{Q}}_4)$ and $\frac{1}{\sqrt{8}}\Lambda_4(\widehat{\mathcal{Q}}_8)$ are isometric for* $p = 47$, *71, 79, 103, 127.*

The proof is analogous to the proof of Theorem 1 and is omitted.

COROLLARY 1. *For $p = 47$ the lattice $\frac{1}{\sqrt{8}}\Lambda(\widehat{\mathcal{Q}}_8)$ has norm* 6, *and the code $\widehat{\mathcal{Q}}_8$ has euclidean minimum weight* 48.

Proof. Follows from the preceding theorem by noticing that $\widehat{\mathcal{Q}}_4$ has euclidean minimum weight 24 [1, 11, 5]. □

The lattice $L_4(\widehat{\mathcal{Q}}_4)$ was considered in [3] and is isometric to P_{48q}. Adopting the definition of P_{48q} in §7.7 of [6], the orthogonal matrix

$$\frac{1}{\sqrt{96}}\begin{pmatrix} -7 & 1 \cdots 1 \\ -1 & \\ \vdots & W - 7I \\ -1 & \end{pmatrix}$$

takes P_{48q} to $L_4(\widehat{\mathcal{N}_4})$ (which is isometric to $L_4(\widehat{\mathcal{Q}}_4)$) by a similar argument to Theorem 1. Similarly it is tantamount to conjecture that the conjectural extremal type II lattice of dimension 80 of example 3 of [13] is taken by

$$\frac{1}{\sqrt{128}}\begin{pmatrix} -7 & 1 \cdots 1 \\ -1 & \\ \vdots & W - 7I \\ -1 & \end{pmatrix}$$

into $L_4(\widehat{\mathcal{N}_4})$.

5. Conclusion

It would be interesting to lift the remaining three Conway-Pless codes over Z_4 and obtain by construction A_4 the three remaining zero-defect lattices of the Koch-Venkov classification. Similarly the construction of P_{48q} by construction B_3 applied to ternary QR codes and density doubling [6, p.149] suggests a construction modulo 6. Eventually, quaternary double circulant codes which produce an even extremal unimodular lattice in dimension 40 [5] should be amenable to a similar analysis.

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