

ANNALES DE L'I. H. P., SECTION A

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Annales de l'I. H. P., section A, tome 48, n° 4 (1988), p. 325-332

http://www.numdam.org/item?id=AIHPA_1988__48_4_325_0

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Uncertainty relations and state spaces

by

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ABSTRACT. — We show that on a quantum logic L which has a sufficient set of states $S(L)$ with the property: for every two noncompatible elements a, b of L there is a state $s \in S(L)$ such that $s(a) = s(b) = 1$, the uncertainty relations cannot be satisfied for any pair of observables on L .

RÉSUMÉ. — Nous montrons que si une logique quantique L a un ensemble d'états $S(L)$ assez grand, c'est-à-dire si pour toute paire a, b d'éléments non compatibles de L il existe un état $s \in S(L)$ tel que $s(a) = s(b) = 1$, alors la relation d'incertitude ne peut être satisfaite pour aucune paire d'observables de L .

1. INTRODUCTION

A quantum logic (a logic in short) is a partially ordered set L with the first and last elements 0 and 1, respectively, and with the orthocomplementation $' : L \rightarrow L$ such that

- i) $(a')' = a$,
- ii) $a \leq b \Rightarrow b' \leq a'$,
- iii) $a \vee a' = 1$,
- iv) for any sequence $\{a_i\} \subset L$ such that $a_i \leq a'_j$ ($i \neq j, i, j = 1, 2, \dots$)

the supremum $\bigvee_{i=1}^{\infty} a_i$ exists in L ,

- v) if $a \leq b$ then there is $c \in L$ such that $c \leq a'$ and $b = a \vee c$.

Two elements $a, b \in L$ are said to be orthogonal (written $a \perp b$) if $a \leq b'$, and $a, b \in L$ are said to be compatible (written $a \leftrightarrow b$) if there are mutually orthogonal elements a_1, b_1, c in L such that $a = a_1 \vee c$, $b = b_1 \vee c$. We have $a \leq b \Rightarrow a \leftrightarrow b$, $a \leftrightarrow b \Rightarrow a \leftrightarrow b'$.

A state on L is a map $s : L \rightarrow [0, 1]$ such that $s(1) = 1$ and $s\left(\bigvee_{i=1}^{\infty} a_i\right) = \sum_{i=1}^{\infty} s(a_i)$ for any sequence $\{a_i\}$ of mutually orthogonal elements of L . Let $S(L)$ denote the set of all states on L , i. e. the state space of L .

A set $S \subset S(L)$ is said to be sufficient if for every $a \in L$, $a \neq 0$, there exists $s \in S$ such that $s(a) = 1$, ordering if $a \not\leq b$ implies that there is $s \in S$ such that $s(a) > s(b)$, strongly ordering if $a \not\leq b$ implies that there is $s \in S$ such that $s(a) = 1$, $s(b) \neq 1$.

A strongly ordering set S is ordering and sufficient, but in general, an ordering and sufficient set of states need not be strongly ordering (see e. g. [1] for the proofs of these statements).

A state $s \in S(L)$ such that $s(a) \in \{0, 1\}$ for all $a \in L$ is called dispersion free or a 0-1 state. Let S_0 be a set of 0-1 states. The conditions— S_0 is ordering—and— S_0 is strongly ordering—are equivalent. Indeed, let S_0 be ordering and let $a \not\leq b$. Then there is $s \in S_0$ such that $s(a) > s(b)$. But this means that $s(a) = 1$ and $s(b) = 0$, i. e. S_0 is strongly ordering.

Let $B(\mathbb{R})$ denote the family of all Borel subsets of the real line \mathbb{R} . An observable on a logic L is a map $x : B(\mathbb{R}) \rightarrow L$ such that

$$i) \quad x(\mathbb{R}) = 1,$$

$$ii) \quad x(E^c) = x(E)' \text{ for any } E \in B(\mathbb{R}), \text{ where } E^c = \mathbb{R} - E,$$

$$iii) \quad x\left(\bigcup_{i=1}^{\infty} E_i\right) = \bigvee_{i=1}^{\infty} x(E_i) \text{ for any sequence } \{E_i\} \text{ of mutually disjoint}$$

elements of $B(\mathbb{R})$.

If x is an observable and $s \in S(L)$, the map $s_x : E \mapsto s(x(E))$ is a probability measure on $B(\mathbb{R})$. The expectation of x in the state s is defined by

$$s(x) = \int t s_x(dt),$$

if the integral on the right exists, and the variance of x in the state s is defined by

$$\text{var}_s(x) = \int (t - s(x))^2 s_x(dt),$$

if the integral on the right exists.

Two observables x, y on L are compatible if $x(E) \leftrightarrow y(F)$ for any $E, F \in B(\mathbb{R})$. The spectrum $\sigma(x)$ of an observable x is the smallest closed

subset C of R such that $x(C) = 1$. An observable x is bounded if $\sigma(x)$ is compact.

We shall need the following lemma.

LEMMA 1. — Let x be an observable on a logic L . Then $t \in \sigma(x)$ if and only if for any open set $U \subset R$ such that $t \in U$ we have $x(U) \neq 0$.

Proof. — Let $t \notin \sigma(x)$. As R is a regular topological space, there are disjoint open sets U, V such that $t \in U$ and $\sigma(x) \subset V$. This implies that $x(U) = 0$. Now let there exist an open set $U \subset R$ such that $t \in U$ and $x(U) = 0$. Then U^c is closed and $x(U^c) = 1$. This implies that $\sigma(x) \subset U^c$, i. e. $t \notin \sigma(x)$.

2. CLASSES OF LOGICS

Let L denote a quantum logic, $S(L)$ the state space of L and $S_0(L)$ the set of all 0-1 states on L . In [6], the following classes of logics with sufficient state spaces have been studied.

$C_1 : a \not\leftrightarrow b \Rightarrow$ there is $s \in S(L)$ such that $s(a) = 1$ and $s(b) \neq 1$,

$C_2 : a \not\leftrightarrow b \Rightarrow$ to any given $\varepsilon > 0$ there is $s \in S(L)$ such that $s(a) = 1$ and $s(b) > 1 - \varepsilon$,

$C_3 : a \not\leftrightarrow b \Rightarrow$ there is $s \in S(L)$ such that $s(a) = s(b) = 1$,

$C_4 : S_0(L)$ is sufficient and $a \not\leftrightarrow b \Rightarrow$ there is $s \in S_0(L)$ such that $s(a) = s(b) = 1$.

Clearly, $C_1 \supset C_2 \supset C_3 \supset C_4$ and by [6], all the inclusions are proper. It is easy to see that C_1 contains exactly logics with strongly ordering state spaces. Indeed, let $S(L)$ be strongly ordering. Since $a \not\leftrightarrow b$ implies $a \not\leq b$, there is $s \in S(L)$ such that $s(a) = 1, s(b) \neq 1$, i. e. $L \in C_1$. On the other hand, let $L \in C_1$ and let $a \not\leq b$. We have only to check the case when $a \leftrightarrow b$. In this case $a = a_1 \vee c, b = b_1 \vee c$, where a_1, b_1, c are mutually orthogonal. The condition $a \not\leq b$ implies that $a_1 \neq 0$. Since $S(L)$ is sufficient, there is $s \in S(L)$ such that $s(a_1) = 1$. This implies that $s(a) = 1$ and $s(b) = 0$, hence $S(L)$ is strongly ordering.

Let H be a Hilbert space. Let $L(H)$ denote the quantum logic of all closed subspaces of H (or equivalently, of all projections on H). The logic $L(H)$ is called a Hilbert space logic. For $M \in L(H)$, let P^M denote the corresponding projection. For any $f \in H, \|f\| = 1$, the map $s_f : M \rightarrow \langle P^M f, f \rangle$, where $\langle \cdot, \cdot \rangle$ is the inner product in H , defines a state on $L(H)$, which is called a vector state. According to Gleason theorem, if $\dim H \geq 3$ and H is separable, every state on $L(H)$ is a σ -convex combination of vector states. Let $M, N \in L(H)$ and let $M \not\leq N$. Then there exists a unit vector $f \in M$, such that $f \notin N$, therefore $s_f(M) = 1, s_f(N) \neq 1$. Hence $L(H)$ belongs to C_1 . Let for any unit vector $f \in H, [f]$ denote the one-dimensional subspace

generated by f . As the only state on $L(H)$ which maps $[f]$ to 1 is s_f , $L(H) \notin C_2$. There has been shown in [7], that the logics of the class C_2 have the following interesting property: for any two bounded observables x, y , the condition $s(x) = s(y)$ in every state $s \in S(L)$ implies that $x = y$. In other words, the logics in C_2 satisfy the condition U (= Uniqueness, see [2], p. 55). The Hilbert space logics also satisfy the property U. In general, it is not known if the logics in C_1 satisfy this condition.

A special family of logics is formed by σ -classes. A σ -class is a family of subsets of a nonempty set X which contains X and is closed under the formations of set-theoretical complements and countable unions of pairwise disjoint elements. A σ -class ordered by inclusion and orthocomplemented by set-theoretical complementation is a quantum logic. By [2], p. 69, a σ -class can be characterized as a logic possessing an ordering set of 0-1 states. It is easy to see that the class C_4 consists exactly of all σ -classes. Indeed, let L be a σ -class. Since the set of all 0-1 states on L is ordering, it is also strongly ordering. Let $a, b \in L$ be such that $a \not\leq b$, then surely $a \not\leq b'$, and therefore there is $s \in S_0(L)$ such that $s(a) = 1$, $s(b') = 0$, i. e. $s(b) = 1$. Hence $L \in C_4$. On the other hand, if $L \in C_4$ then using similar arguments to that used by proving that a logic $L \in C_1$ has a strongly ordering state space, we show that L is a σ -class.

Let H be a two-dimensional Hilbert space. Then every set of non-zero mutually orthogonal elements in $L(H)$ is of the form $\{a, a'\}$, $a \in L(H)$. It is easy to see that $L(H)$ is a σ -class. Indeed, let

$$S_0 = \{s : L(H) \rightarrow \{0, 1\} \mid s(a) + s(a') = 1\}$$

and $h(a) = \{s \in S_0 \mid s(a) = 1\}$. It is easy to check that the mappings in S_0 are states on $L(H)$, S_0 is ordering and the family $\Delta = \{h(a) \mid a \in L(H)\}$ of subsets of S_0 forms a σ -class. To give a more explicit representation, let $H = \mathbb{R}^2$ and let $X = [0, \pi) \times [0, \pi)$. Put

$$\tau(\alpha) = \begin{cases} [0, \alpha) \times \left[\alpha + \frac{\pi}{2}, \pi\right) & \text{if } 0 \leq \alpha < \frac{\pi}{2} \\ \left[\alpha - \frac{\pi}{2}, \pi\right) \times [0, \alpha) & \text{if } \frac{\pi}{2} \leq \alpha < \pi. \end{cases}$$

It is easy to check that $\Delta = \{\emptyset, X, \tau(\alpha) \mid \alpha \in [0, \pi)\}$ is a σ -class ($\tau(\alpha) \cap \tau(\beta) = \emptyset$ iff $\beta = \alpha + \frac{\pi}{2}$, $\tau(\alpha)^c = \tau\left(\alpha + \frac{\pi}{2}\right)$). Every one-dimensional subspace in \mathbb{R}^2

can be characterized by an angle α , $\alpha \in [0, \pi)$. Denote by $[\alpha]$ the one-dimensional subspace corresponding to α . The map $h : L(H) \rightarrow \Delta$, $h(0) = \emptyset$, $h(H) = X$, $h([\alpha]) = \tau(\alpha)$, defines an isomorphism between $L(H)$ and Δ . Let $s_{(\beta, \gamma)}$ be the probability measure on Δ concentrated in the point $(\beta, \gamma) \in X$. The set $S_0 = \{s_{(\beta, \gamma)} \mid (\beta, \gamma) \in X\}$ represents the set of all 0-1 states on $L(H)$.

3. UNCERTAINTY RELATIONS

Let x be an observable on a logic L . We put

$$V(x) = \{ s \in S(L) \mid \text{var}_s(x) < \infty \}.$$

For any two observables x, y on L , one of the following alternative possibilities occurs:

(A) $(\forall \varepsilon > 0)(\exists s \in V(x) \cap V(y) (\text{var}_s(x) \cdot \text{var}_s(y) < \varepsilon)$

(B) $(\exists \varepsilon > 0)(\forall s \in V(x) \cap V(y) (\text{var}_s(x) \cdot \text{var}_s(y) \geq \varepsilon).$

If (B) occurs, we say that the uncertainty relation holds for x and y (see [3], [1]).

For $t \in R, \delta > 0$, put $U(t, \delta) = \{ r \in R \mid |t - r| < \delta \}$. If x is an observable and $t \in \sigma(x)$, then $x(U(t, \delta)) \neq 0$ by Lemma 1.

Let x and y be observables. The following two possibilities can occur:

(a) $(\forall(u, v, \delta) : u \in \sigma(x), v \in \sigma(y), \delta > 0) (\exists \eta_0 > 0) (\forall \eta, 0 < \eta < \eta_0) (x(U(u, \delta)) \leftrightarrow y(U(v, \eta))).$

(b) $(\exists(u, v, \delta) : u \in \sigma(x), v \in \sigma(y), \delta > 0) (\forall \eta_0 > 0) (\exists \eta, 0 < \eta < \eta_0) (x(U(u, \delta)) \not\leftrightarrow y(U(v, \eta))).$

THEOREM 1. — Let L be a logic with a sufficient state space. If for the observables x and y on L the condition (a) holds, then the uncertainty relation does not hold. In other words, (a) \Rightarrow (A).

Proof. — Let (a) hold for the observables x and y . We show that the following holds:

$$(\forall(u, \delta) : u \in \sigma(x), \delta > 0) (\exists v \in \sigma(y) (\forall \eta < \eta_0) (x(U(u, \delta)) \wedge y(U(v, \eta)) \neq 0)$$

(η_0 exists by (a)). Suppose that the opposite holds, i. e.

$$(\exists(u, \delta) : u \in \sigma(x), \delta > 0) (\forall v \in \sigma(y)) (\exists \eta < \eta_0) (x(U(u, \delta)) \wedge y(U(v, \eta)) = 0).$$

Since by (a) $x(U(u, \delta)) \leftrightarrow y(U(v, \eta))$, it is $x(U(u, \delta)) \perp y(U(v, \eta))$. We have $\sigma(y) \subset \cup \{ U(v, \eta(v)) \mid v \in \sigma(y) \}$. By the second countability of the topology of R , there is a countable set $\{ v_i \}$ such that

$$\sigma(y) \subset \bigcup_{i=1}^{\infty} U(v_i, \eta_i) \quad \text{and}$$

$$y\left(\bigcup_{i=1}^{\infty} U(v_i, \eta_i)\right) = \bigvee_{i=1}^{\infty} y(U(v_i, \eta_i)) \geq y(\sigma(y)) = 1.$$

Then $x(U(u, \delta)) \leq y(U(v_i, \eta_i))'$ for all $i = 1, 2, \dots$ implies that

$$x(U(u, \delta)) \leq \bigwedge_{i=1}^{\infty} y(U(v_i, \eta_i))' = \left(\bigvee_{i=1}^{\infty} y(U(v_i, \eta_i)) \right)' = 0,$$

which contradicts the supposition that $u \in \sigma(x)$. Let us choose $u \in \sigma(x)$ and $\delta > 0$. Then there is $v \in \sigma(y)$ such that for any $\eta < \eta_0$ ($\eta_0 = \eta_0(u, v, \delta)$) we have $x(U(u, \delta)) \wedge y(U(v, \eta)) \neq 0$. By the sufficiency of $S(L)$ there is $s \in S(L)$ such that

$$s(x(U(u, \delta)) \wedge y(U(v, \eta))) = 1.$$

Hence $\text{var}_s(x) = \int_{U(u, \delta)} (t - s(x))^2 s_x(dt) < 4\delta^2$, and similarly $\text{var}_s(y) < 4\eta^2$.

By choosing $\eta < \min\left(\eta_0, \frac{\sqrt{\varepsilon}}{2\delta}\right)$, we obtain that for any given $\varepsilon > 0$ there exists a state $s \in S(L)$ such that $\text{var}_s(x) \cdot \text{var}_s(y) < \varepsilon$.

Remark. — Condition (a) can be weakened to (a'), where

$$(a') (\exists(u, \delta) : u \in \sigma(x), \delta > 0) (\forall v \in \sigma(y)) (\exists \eta_0 > 0) (\forall \eta, 0 < \eta < \eta_0) (x(U(u, \delta)) \leftrightarrow y(U(v, \eta)))$$

and (a') \Rightarrow (A).

THEOREM 2. — Let L be a logic which is a lattice. If for the observables x and y the condition (a) holds, then x and y are compatible.

Proof. — Let U be any open subset of R . We have $y(U) = y(U \cap \sigma(y))$ and $U \cap \sigma(y) \subset \cup \{U(v, \eta(v)) \mid v \in \sigma(y) \cap U\} \subset U$, where $\eta(v) > 0$. By the second countability of R , there is a countable subfamily $\{U(v_i, \eta(v_i))\}$ such that

$$U \cap \sigma(y) \subset \bigcup_{i=1}^{\infty} U(v_i, \eta(v_i))$$

and

$$y(U) = \bigvee_{i=1}^{\infty} y(U(v_i, \eta(v_i))).$$

By the property (a), to any $u \in \sigma(x)$ and $\delta > 0$, and to any open set U there are $v_i \in \sigma(y)$, $\eta(v_i) > 0$ such that $y(U) = \bigvee_{i=1}^{\infty} y(U(v_i, \eta(v_i)))$ and

$$x(U(u, \delta)) \leftrightarrow y(U(v_i, \eta(v_i)))$$

for $i = 1, 2, \dots$, which implies that $x(U(u, \delta)) \leftrightarrow y(U)$. Now let V be an open subset of R . Then there are $u_i \in \sigma(x)$ and $\delta_i > 0$, $i = 1, 2, \dots$ such

that $x(V) = \bigvee_{i=1}^{\infty} x(U(u_i, \delta_i))$. Since $x(U(u_i, \delta_i)) \leftrightarrow y(U)$, we get $x(V) \leftrightarrow y(U)$

for any open subsets U, V of \mathbb{R} , and this implies that $x \leftrightarrow y$.

THEOREM 3. — Let $L \in C_3$. Then the uncertainty relation (B) does not hold for any pair of observables on L .

Proof. — Let x, y be observables on L . By Theorem 1, (a) \Rightarrow (A). Suppose that (b) holds for x and y . Then there are $u \in \sigma(x), \delta > 0, v \in \sigma(y)$ such that to any $\eta_0 > 0$ there is $\eta < \eta_0$ such that $x(U(u, \delta)) \not\leftrightarrow y(U(v, \eta))$. As L belongs to C_3 , there is a state $s \in S(L)$ such that $s(x(U(u, \delta))) = s(y(U(v, \eta))) = 1$. Choosing η sufficiently small we obtain $\text{var}_s(x) \cdot \text{var}_s(y) < \varepsilon$ for any given $\varepsilon > 0$.

Let $H^2 = L(\mathbb{R})$ be the set of all square-integrable complex valued functions defined on \mathbb{R} with respect to the Lebesgue measure. Let q and p be the « position » and « momentum » observables corresponding to the self-adjoint operators P, Q , where $(Qf)(r) = rf(r), (Pf)(r) = -ih \frac{d}{dr} f(r)$

for $r \in \mathbb{R}$. It can be shown that $\text{var}_{s_f}(q) \cdot \text{var}_{s_f}(p) \geq \frac{h^2}{4}$ for all $f \in D(Q) \cap D(P)$,

where $D(A)$ denotes the domain of the operator A (see e. g. [8], p. 77, 393-394 for the proof). For any self-adjoint operator A its domain

$$D(A) = \left\{ f \in H \mid \int t^2 \langle E^A(dt)f, f \rangle < \infty \right\} = \{ f \in H \mid s_f \in V(E^A) \},$$

where E^A is the spectral measure (which can be identified with the observable corresponding to A by the spectral theorem). Owing to Gleason theorem, every state on $L(H)$ is a σ -convex combination of vector states. From this we may conclude that the observables p and q satisfy the uncertainty relation in the sense of our definition.

The above example shows that there are couples of observables on the logics of the class C_1 which satisfy the uncertainty relations. It remains an open question if there exist couples of observables on the logics of the class C_2 satisfying the uncertainty relations.

In [3], the notion of complementarity has been introduced as follows. Let x, y be observables on a logic L . We say that x, y are complementary if $x(E) \wedge y(F) = 0$ for every bounded subsets E, F of \mathbb{R} such that $x(E) \neq 1$ and $y(F) \neq 1$. It is a well-known fact that the observables q, p in the above example are complementary (see e. g. [4], [5]). Now let us consider the logic $L(H)$ of the two-dimensional Hilbert space H . It is easy to see that any two noncompatible observables on $L(H)$ are complementary. This

example shows that complementarity is not excluded on the logics of the class C_3 or even C_4 . However, it would be interesting to find less trivial examples of unbounded complementary observables on the logics of the class C_3 .

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(Manuscrit reçu le 3 décembre 1987)

(Version révisée reçue le 9 mars 1988)