



## Logic

# Claws in digraphs

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### Abstract

We study in this Note the existence of claws in digraphs. We extend a result of Saks and Sós to the tournament-like digraphs. *To cite this article: A. El Sahili, M. Kouider, C. R. Acad. Sci. Paris, Ser. I 346 (2008).*

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### Résumé

**Arbres simples dans les graphes orientés.** Nous étudions dans cette Note le problème de l'existence des arbres simples dans les graphes orientés. Nous étendons un résultat de Saks et Sós aux graphes orientés presque complets. *Pour citer cet article : A. El Sahili, M. Kouider, C. R. Acad. Sci. Paris, Ser. I 346 (2008).*

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## 1. Introduction

The digraphs considered here are all connected, they have no loops, multiple edges or circuits of length two. The chromatic number of a digraph is the chromatic number of its underlying graph. A *block* of a path in a digraph is a maximal directed subpath. We recall that the *length* of a path is the number of its edges.

A *rooted tree* is a tree  $T$  with a specified vertex, called the *root* of  $T$ . A *branching* (resp. *inbranching*) is an orientation of a rooted tree in which every vertex except the root has in degree 1 (resp. out degree 1). The *level* of a vertex in a branching is its distance from the root. A *claw* is an inbranching in which only the root may have an in degree more than 2. It will be denoted by  $P(s_1, s_2, \dots, s_k)$  if it is formed by  $k$  directed paths ending at the root, of lengths  $s_1, s_2, \dots, s_k$  respectively. In particular,  $P(k, l)$  is a path with two blocks of lengths  $k$  and  $l$ .

Saks and Sós proved the following result about claws in tournaments:

**Theorem 1.** (See [5].) *Any  $2(n - 1)$ -tournament contains any claw of order  $n$ .*

This result became a corollary of that one proved later by Havet and Thomassé [3] showing that a  $2(n - 1)$ -tournament contains any branching of order  $n$ .

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By forbidding some paths, we proved [1] the following:

**Theorem 2.** *Let  $D$  be an  $2(n - 1)$ -chromatic digraph without paths  $P(2n - 3, 1)$ . Then  $D$  contains any claw of order  $n$ .*

This gives a direct proof of Saks and Sós' theorem by remarking that the path  $P(2n - 3, 1)$  contains  $2n - 1$  vertices.

In order to generalize the Saks and Sós' result to the arbitrary digraphs, we want in this paper to extend it to some type of digraphs, the tournament-like digraphs (defined below). We treat in particular the case when the claw is a path with two blocks.

## 2. Spanning branching forest

A *branching forest* is a digraph each of whose components is a branching. The level of a vertex  $v$  in a branching forest  $F$ , denoted by  $l_F(v)$ , is the order of the directed path in  $F$  ending at  $v$  from the root of the branching of  $F$  containing  $v$ . The set  $L_i(F)$  denotes the set of all vertices of level  $i$  in  $D$ . We denote by  $l(F)$  the maximal integer  $i$  such that  $L_i(F)$  is non-empty. Again if  $v$  is a vertex in  $F$ , we denote by  $T_v(F)$  the branching of  $F$  having  $v$  as root.

**Definition 1.** A spanning branching forest  $F$  of a digraph  $D$  is said to be maximal if the sum  $\sum_{v \in D} l_F(v)$  is maximal.

**Definition 2.** A tournament-like digraph  $D$  is a digraph containing a maximal spanning branching forest  $F$  such that  $l(F) = \chi(D)$ .

An example of a tournament-like digraph which is not a tournament can be obtained easily from any  $n$ -chromatic simple graph  $G$ : Put  $V(G) = \bigcup_{i=1}^n S_i$  where  $S_i$  are stable sets in  $G$ . An edge  $x_i x_j \in E(G)$ ,  $x_i \in S_i$ ,  $x_j \in S_j$  is oriented from  $x_i$  to  $x_j$  whenever  $i < j$ . The obtained acyclic digraph has no directed path of length  $n$  so it is a tournament-like digraph. An example with circuit (see [2]) can be formed by an odd oriented cycle  $C$  containing no directed paths of length 3, such that any vertex in  $C$  is joined to all the vertices of a strongly connected tournament on  $n - 3$  vertices.

**Lemma 1.** *Let  $F$  be a maximal spanning branching forest of a digraph  $D$ . Then the sets  $L_i(F)$ ,  $i \geq 0$ , are stable sets in  $D$ .*

By simply remarking that any vertex in  $L_i(F)$  is the terminus of a unique directed path of length  $i - 1$  in  $F$ , the Roy–Gallai theorem [2–4] will be a simple consequence of Lemma 1.

Lemma 1 is also used in [1] to prove the following theorem:

**Theorem 3.** *Every  $n + 1$ -chromatic digraph contains any path with two blocs of length  $n - 1$ .*

The following property of a maximal spanning branching forest of a digraph  $D$  will be useful in the sequel:

**Lemma 2.** *Let  $F$  be a maximal spanning branching forest of a digraph  $D$ . Let  $e = (x, y) \in E(D)$  with  $l_F(x) > l_F(y)$ , then  $F$  contains a directed  $yx$ -path.*

Consequently, we have the following property:

**Lemma 3.** *Let  $F$  be a maximal spanning branching forest of a digraph  $D$ . Set  $l(F) = l$ . For every  $y \in L_l(F)$  and for every  $r < l$ , we have either  $N(y) \cap L_r(F) = N^+(y) \cap L_r(F)$  or  $N(y) \cap L_r(F) = N^-(y) \cap L_r(F)$ .*

We define now an important operation on maximal forests:

**Lemma 4.** *Let  $F$  be a maximal spanning branching forest of a digraph  $D$  with  $l(F) = l$ . Let  $C = v_r v_{r+1} \cdots v_{l-1} v_l$  be a circuit of  $D$  such that  $v_i \in L_i(F)$ ,  $1 < r \leq i \leq l$ . Suppose that there is a vertex  $x \in L_{r-1}(F)$  such that  $e = (x, v_l) \in E(D)$ . Then the forest  $F'$  obtained from  $F$  by rotating the cycle  $C$  in the negative sense, that is  $F' =$*

$F + e + (v_l, v_r) - (v_{l-1}, v_l) - f$ , where  $f$  denotes the arc of  $F$  with head  $v_r$ , is a maximal spanning branching forest and so  $T_{v_r}(F) = v_r v_{r+1} \cdots v_{l-1} v_l$  and  $l(F') = l$ .

$F'$  is called a rotating of  $F$ , it is denoted by  $\mathfrak{R}(F)$  or  $\mathfrak{R}_{C,x}(F)$  for more precision.

Note that the above operation can be also defined even if  $r = 1$  by simply rotating  $C$ , that is  $F' = F + (v_l, v_r) - (v_{l-1}, v_l)$ .

Consider a maximal spanning branching forest  $F$  of a digraph  $D$  with  $l(F) = l$ . A vertex  $y \in L_l(F)$  is said to have a root  $r$ , we write  $r(y) = r$  if  $N^+(y) \cap L_r(F) \neq \emptyset$  and  $N^-(y) \cap L_i(F) \neq \emptyset$  for all  $i \leq r - 1$ , if any. Set  $U_r(F) = \{y \in L_l(F), r(y) = r\}$ ,  $R(D) = \bigcup_{F,r} U_r(F)$  and  $\mathfrak{R}(D) = \{r(y), y \in R(D)\}$ . All these sets may be empty.

**Lemma 5.** *Let  $F$  be a maximal spanning branching forest of a digraph  $D$  with  $l(F) = l$ . Let  $u \in U_r(F)$  for some  $r < l$ , and let  $C = v_r v_{r+1} \cdots v_{l-1} u$  be the circuit of  $D$  such that  $v_i \in L_i(F)$ ,  $r \leq i \leq l - 1$ . If  $F' = \mathfrak{R}(F)$  for some rotating of  $F$  using  $C$ , then we have:  $U_t(F) = U_t(F')$  for all  $t < r$  and either  $U_r(F') = U_r(F) - \{u\}$  or  $U_r(F') = (U_r(F) - \{u\}) \cup \{v_{l-1}\}$ .*

If  $\mathfrak{R}(D) \neq \emptyset$ , set  $\mathfrak{R}(D) = \{r_0, r_1, \dots, r_s\}$  with  $r_0 < r_1 < \dots < r_s$ ,  $\mathcal{F} = \{F \text{ maximal, } l(F) \text{ minimal}\}$  and define the following decreasing sequence of non-empty classes of maximal forests:  $\mathcal{F}_0 = \{F \in \mathcal{F}, |U_{r_0}(F)| \text{ is minimal}\}$  and  $\mathcal{F}_{i+1} = \{F \in \mathcal{F}_i, |U_{r_{i+1}}(F)| \text{ is minimal}\}$ ,  $1 \leq i < s$ . Based on the above lemma we may remark easily that  $\mathcal{F}_k$ ,  $1 \leq k \leq s$  is closed by some one of the rotating operations:

**Lemma 6.** *Let  $F \in \mathcal{F}_k$ ,  $1 \leq k \leq s$ . Let  $u$  be a vertex in  $U_{r_k}(F)$ , let  $C = v_{r_k} v_{r_k+1} \cdots v_{l-1} u$  be the circuit of  $D$  such that  $v_i \in L_i(F)$ ,  $r_k \leq i \leq l - 1$ . If  $F' = \mathfrak{R}(F)$  for some rotating of  $F$  using  $C$ , then  $F' \in \mathcal{F}_k$  and  $U_{r_k}(F') = (U_{r_k}(F) - \{u\}) \cup \{v_{l-1}\}$ .*

Forests in  $\mathcal{F}_k$  have the following crucial property:

**Lemma 7.** *Let  $F \in \mathcal{F}_s$ . Let  $u$  be a vertex in  $U_{r_k}(F)$ ,  $1 \leq k \leq s$ , let  $C = v_{r_k} v_{r_k+1} \cdots v_{l-1} u$  be the circuit of  $D$  such that  $v_i \in L_i(F)$ ,  $r_k \leq i \leq l - 1$ . Then for all  $v \in C$  we have  $N^-(v) \cap L_j(F) \neq \emptyset$  for  $j < r_k$ , and  $N(v) \cap (L_j(F) - C) = \emptyset$  for  $j \geq r_k$ .*

### 3. Claws

Consider a tournament-like digraph  $D$  and let  $F$  be a maximal spanning branching forest  $F$  such that  $l(F) = \chi(D)$ . Tournament-like digraphs have the following interesting property:

**Lemma 8.** *A tournament-like digraph  $D$  has a maximal forest  $F$  and a tournament  $T$  with  $V(T) = \{v_r, v_{r+1}, \dots, v_{l-1}, v_l\}$  such that  $v_i \in L_i(F)$ ,  $N^-(v_i) \cap L_j(F) \neq \emptyset$   $r \leq i \leq l$  and  $j < r$ .*

**Proof.** Let  $F \in \mathcal{F}_s$ . We have  $l = l(F) = \chi(D)$ . We suppose that  $L_l(F)$  contains no vertex  $v$  such that  $N^-(v) \cap L_j(F) \neq \emptyset$  for all  $j < l$  since otherwise we take  $V(T) = \{v\}$ . There is a  $k \in \{1, 2, \dots, s\}$  such that  $U_{r_k}(F) \neq \emptyset$  since otherwise  $L_l(F) \cap R(D) = \emptyset$  and so, due to the first supposition, for all  $u \in L_l(F)$  there exists  $j < l$  such that  $N(u) \cap L_j(F) = \emptyset$ . Hence the vertices in  $L_l(F)$  may be distributed over the other levels and as  $l(F) = \chi(D)$ , we obtain a  $(\chi(D) - 1)$ -coloring. A contradiction. Suppose to the contrary that the lemma does not hold. Consider a vertex  $u \in U_{r_k}(F)$  and let  $C = v_{r_k} v_{r_k+1} \cdots v_{l-1} v_l = u$  be the circuit of  $D$  such that  $v_i \in L_i(F)$ ,  $r_k \leq i \leq l$ . Since  $D(V(C))$  is not a tournament then there exist  $i$  and  $j$ ,  $r_k \leq i, j \leq l$ , such that  $v_i v_j \notin E(G(D))$ . By defining the periodic sequence of maximal forests in  $\mathcal{F}_s$  as in Lemma 7, we may replace  $F$  by  $F_{l-i}$  which is always in  $\mathcal{F}_s$ . The vertex  $u$  is replaced by  $v_i$ . Now  $r = l_{F_{l-i}}(v_j) \geq r_k$ . By Lemma 7,  $N(v_i) \cap (L_r(F_{l-i}) - \{v_j\}) = \emptyset$ . Then  $N(v_i) \cap L_r(F_{l-i}) = \emptyset$  as  $v_i v_j \notin E(G(D))$ . If we proceed similarly for all the vertices in  $\bigcup_{i=1}^s U_{r_i}(F)$ , we get a maximal forest  $F' \in \mathcal{F}_s$  such that for all  $u \in L_l(F')$  there exists  $j < l$  such that  $N(u) \cap L_j(F') = \emptyset$  which leads to a contradiction. This completes the proof of the lemma.  $\square$

We prove now our main result:

**Theorem 4.** *A tournament-like digraph  $D$  with  $\chi(D) = 2n - 2$  contains any claw on  $n$  vertices.*

**Proof.** Let  $C = P(s_1, s_2, \dots, s_k)$  be a claw of order  $n$  ( $s_0 + s_1 + \dots + s_k = n - 1$ ) with  $s_0 = 0$ . Let  $T$  be a tournament on  $t$  vertices and  $F$  be a maximal forest as in the above lemma. Let  $i$ ,  $0 \leq i \leq k$ , be the maximal integer such that  $s_0 + s_1 + \dots + s_i \leq \frac{t}{2}$ . The case  $i = k$  corresponds to the case  $t = 2n - 2$  and so by the above lemma  $D$  is a tournament on  $2n - 2$  vertices, it contains a copy of  $C$  by Theorem 1. Then we suppose that  $i < k$ , we have  $s_0 + s_1 + \dots + s_i + s = \frac{t}{2}$  with  $0 \leq s < s_{i+1}$ . Let  $s'$  and  $r$  be such that  $s + s' = s_{i+1}$  and  $r + t = 2n - 2$ . We have  $s' + s_{i+2} + \dots + s_k = \frac{t}{2} < r$ . Thus  $T$  is a tournament on  $t = 2(s_0 + s_2 + \dots + s_i + s)$  vertices, it contains a copy of the claw  $C' = P(s_0, s_1, \dots, s_i, s)$ . Call  $v_0$ , the root of  $C'$  and let  $v$  be the tail of the leaf corresponding to the path of  $C'$  of length  $s$ . By the above lemma, there is an in neighbor  $u$  of  $v$  in  $L_{r-1}(F)$ . Consider a directed  $w_1 u$ -path  $P_1$  of length  $s' - 1$  in  $F$  where  $w_1 \in L_{r-s'}(F)$ . The path  $P_1$  together with  $w_1 u v \dots v_0$  form the  $(i + 1)$ th branch of  $C$ . Suppose now that the  $(i + j)$ th branch of  $C$  is constructed without using vertices in  $L_h(F)$ ,  $h < a = r - (s' + s_{i+2} + \dots + s_{i+j})$ . if  $i + j < k$  then  $a \geq s_{i+j+1}$  and so if we consider an in neighbor of  $v_0$  in  $L_a(F)$ , we may construct the  $(i + j + 1)$ th branch of  $C$  without using vertices in  $L_h(F)$ ,  $h < a - s_{i+j+1} = r - (s' + s_{i+2} + \dots + s_{i+j+1})$ . This induction completes the construction of a copy of  $C$  in  $D$ .  $\square$

A natural problem is to ask about the characterization of tournament-like digraphs.

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